

Claims

We claim:

- 1 1. A method for decoding a received word for a finite geometry code, comprising:
2 a one-time initialization procedure; comprising:
3 defining a parity check matrix for the finite geometry code;
4 representing each bit in the received word by an element, each
5 element having either a zero, one or uncertain value; and
6 an iterative decoding procedure, comprising:
7 terminating when a termination condition is true, and otherwise:
8 determining a set of votes for each element depending on the parity
9 check matrix and current values of the elements;
10 updating the elements based on the set of votes for each element and
11 the received word.
- 1 2. The method of claim 1, in which the parity check matrix representing an (N, k)
2 finite geometry code has M rows and N columns and rank $N - k$, and the M rows are
3 incidence vectors of a set of selected μ -flats in the finite geometry with respect to
4 N points of the finite geometry, each row containing J ones and $N - J$ zeros.
- 1 3. The method of claim 1, in which the μ -flats are cyclic shifts of each other.
- 1 4. The method of claim 1, in which each of the elements express a state of the
2 decoder corresponding to a bit of the received word, and the termination condition
3 is satisfied when the state of the decoder is a code-word.

1 5. The method of claim 1, in which the initialization procedure and the decoding
2 procedure are restarted with a substantially larger parity check matrix if the
3 termination condition is satisfied and the state of the decoder does not correspond
4 to a code-word.

1 6. The method of claim 1, in which each row of the parity check matrix
2 corresponds to a parity check, and each parity check sends votes to J elements,
3 where the elements that a parity check sends a vote to correspond to the columns of
4 the parity check matrix that have the value 1 in a row corresponding to the parity
5 check.

1 7. The method of claim 6, in which the vote for a particular parity check to a
2 particular element is determined by states of the $J-1$ other elements associated with
3 the parity check.

1 8. The method of claim 6, in which the vote from a particular parity check to a
2 particular element is an abstention if any of the $J-1$ other elements associated with
3 the parity check has the uncertain value, and is one when the number of other
4 elements associated with the parity check having a value of one is odd, and is zero
5 otherwise.

1 9. The method of claim 6, further comprising:
2 determining a recommendation and strength of recommendation for each
3 element from the associated parity checks that send a vote to it.

- 1 10. The method of claim 9, in which the recommendation is zero when a majority
2 of the votes are zero, the recommendation is one when the majority of the votes are
3 one, and there is no recommendation when the zero and one votes are equal.
- 1 11. The method of claim 10, in which the strength of the recommendation is a
2 magnitude of a difference between the number of zero votes and the number of one
3 votes.
- 1 12. The method of claim 11, in which the votes determine a next updated state of a
2 particular element by means of a comparison of the strength of the
3 recommendation of the parity checks with a threshold b_{flip} and a threshold $b_{uncertain}$.
- 1 13. The method of claim 12, in which the next updated state of a particular element
2 is equal to a value in the received word if the recommendation of the parity checks
3 agrees with the value in the received word.
- 1 14. The method of claim 12, in which the next updated state of a particular element
2 is equal to a value in the received word if the recommendation of the parity checks
3 is not equal to the value in the received word, and the strength of the
4 recommendation is less than the threshold $b_{uncertain}$, and the next updated state of
5 the particular element is a state representing uncertainty if the recommendation of
6 the parity checks is not equal to the value in the received word, and the strength of
7 the recommendation is greater than or equal to the threshold $b_{uncertain}$, and less than
8 the threshold b_{flip} , and the next updated state of the particular element agree with
9 the recommendation of the parity checks if the strength of the recommendation is
10 greater than the threshold b_{flip} .

- 1 15. The method of claim 12 further comprising:
- 2 updating the values of the thresholds at an end of each decoding cycle,
- 3 according to a pre-determined schedule.